



## Laceability in Hanoi Graphs

R. A. Daisy Singh<sup>1,\*</sup> and R. Murali<sup>2</sup>

<sup>1</sup>B. N. M. Institute of Technology, Bengaluru, Karnataka, India

<sup>2</sup>Dr. Ambedkar Institute of Technology, Bengaluru, Karnataka, India

\*Corresponding author: [thavanya2010@ymail.com](mailto:thavanya2010@ymail.com)

**Abstract.** The topological structure of an interconnection network can be modelled by a connected, simple and undirected graph  $G = (V, E)$  where  $V$  represents the set of processors and  $E$  represents the set of communication links. Interconnection networks are used to interconnect the processors of data centres and cluster computers. The study of Hamiltonicity and the related areas such as Hamiltonian laceability and Hamiltonian connectedness has lot of significance in computer networks. A network (graph) is Hamiltonian connected if it contains a Hamiltonian path between two distinct nodes (vertices). In this paper we shall study the laceability properties associated with Hanoi graphs  $H_n$ . To be more specific we shall explore Hamiltonian- $t^*$ -connectedness of Hanoi graphs  $H_n$  for  $n \geq 3$ .

**Keywords.** Hamiltonian Laceability; Hamiltonian connected

**MSC.** 05C45

**Received:** January 20, 2018

**Accepted:** October 25, 2018

Copyright © 2019 R. A. Daisy Singh and R. Murali. *This is an open access article distributed under the Creative Commons Attribution License, which permits unrestricted use, distribution, and reproduction in any medium, provided the original work is properly cited.*

### 1. Introduction

The Hanoi graph  $H_n$  corresponds to the allowed moves in the tower of Hanoi problem. The tower of Hanoi puzzle invented in 1883 by the French mathematician Edouard Lucas, has become a classic example in the analysis of algorithms and discrete mathematical structures. The Hanoi graph can be constructed by taking the vertices to be the odd binomial coefficients of Pascal's triangle computed on the integers from 0 to  $2^n - 1$  and drawing an edge whenever coefficients are adjacent diagonally or horizontally. The graph  $H_n$  has  $3n$  vertices and  $\frac{3(3^n)-1}{2}$  edges. The diameter of  $H_n$  is  $2^n - 1$ . Each Hanoi graph has exactly two distinct directed

Hamiltonian cycles.  $H_n$  has  $3^n - 1$  small triangles, each of which can contain at most one vertex in an independent vertex set. But the triangles are arranged in a plane in such a way that choosing the apex of each, gives a maximum independent vertex set. Hanoi graphs are perfect. Klavzar and Milutinovic [15], and Aumann et al. have explored [14], some classical numbers of Hanoi graphs. Xuemiao [12] has introduced the notion of Hanoi graph for the towers of Hanoi puzzle which greatly helps investigate the problem. In [13] Kumar and Maheshwari have discussed about the Matrix representation of Hanoi graphs. The tower of Hanoi problem on path graphs have been discussed by Berend et al. [5]. Also, they have discussed about the diameter of Hanoi graphs. Hamiltonian- $t$ -laceability in the brick product of even cycles was studied by Shenoy and Murali [11]. In [7] Girisha and Murali have studied Hamiltonian- $t^*$ -laceability of 4 regular graphs.

A connected graph  $G$  is termed as Hamiltonian  $t_0^*$  ( $t_e^*$ ) connected if there exists in  $G$ , a Hamiltonian path between atleast one pair of its vertices  $u$  and  $v$  with the property  $d(u,v) = t$  for all odd (even)  $t$  such that  $1 \leq t \leq \text{diam}(G)$  where  $\text{diam}(G)$  is the diameter of a graph  $G$ . A connected graph  $G$  is termed as Hamiltonian- $t^*$ -connected if there exists in  $G$ , a Hamiltonian path between atleast one pair of its vertices  $u$  and  $v$  with the property  $d(u,v) = t$  for all  $t$  such that  $1 \leq t \leq \text{diam}(G)$ . The Hanoi graph  $H_3$  is illustrated in Figure 1.

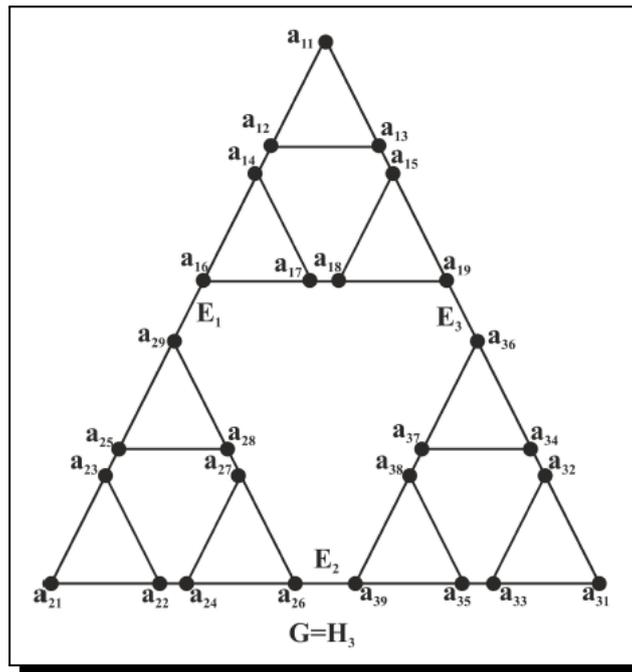


Figure 1

**Definition 1.1.** Let  $P_\alpha$  and  $P_\beta$  respectively be the paths from  $a_i$  to  $a_j$  and  $a_j$  to  $a_k$  in  $G$ . Then the path  $P_\alpha \cup P_\beta$  is the path obtained by extending the path  $P_\alpha$  from  $a_i$  to  $a_j$  and from  $a_j$  to  $a_k$  through the common vertex  $a_j$ . That is if  $P_\alpha : a_i \rightarrow a_j$  and  $P_\beta : a_j \rightarrow a_k$  then  $P_\alpha \cup P_\beta : a_i \rightarrow a_j \rightarrow a_k$ .

## 2. Results

**Lemma 2.1.** *The graph  $H_3$  is  $(t_o^*)$  connected for odd  $t$  such that  $1 \leq t \leq \text{diam}G$ .*

*Proof.* Let  $G = H_n, n = 3$ . Clearly,  $\text{diam}(G) = 2^3 - 1 = 7$ . Let the vertex set of  $H_3$  be  $V(H_3) = \{a_{ij}; i = 1; 1 \leq j \leq 3^{n-1}\} \cup \{a_{ij}; i = 2; 1 \leq j \leq 3^{n-1}\} \cup \{a_{ij}; i = 3; 1 \leq j \leq 3^{n-1}\}$  and the edge set of  $H_3$  be  $E(H_3) = \left\{b_i, 1 \leq i \leq \frac{3^{n+1}-9}{6}\right\} \cup \left\{b'_i, 1 \leq i \leq \frac{3^{n+1}-9}{6}\right\} \cup \left\{b''_i, 1 \leq i \leq \frac{3^{n+1}-9}{6}\right\} \cup \{E_3\}$  where  $\{E_3\} = \{e_1, e_2, e_3\}$  such that  $e_1 = (a_{(i)(j-m)} - a_{(i+1)(j)})$ ,  $e_2 = (a_{(i+1)(j-m)} - a_{(i+2)(j)})$ ,  $e_3 = (a_{(i+2)(j-m)} - a_{(i)(j)})$  and  $m = 2^{(n-1)} - 1$ .

Now to establish the result we consider the following cases.

In each case, we take  $i = 1, j = 3^{n-1}$ .

**Case (i):** Let  $t = 1$ . In  $G, d(a_{12}, a_{11}) = 1$  and the Hamiltonian path between the two vertices  $a_{12}$  and  $a_{11}$  is  $P_1 : \{P_1 \cup P_2 \cup P_3 \cup P_4\} \cup \{E_3\}$ , where  $P_1 : a_{(i)(j-7)} - a_{(i)(j-5)} - a_{(i)(j-2)} - a_{(i)(j-3)}$ ,  $P_2 : a_{(i+1)(j)} - a_{(i+1)(j-1)} - a_{(i+1)(j-4)} - a_{(i+1)(j-6)} - a_{(i+1)(j-8)} - a_{(i+1)(j-7)} - a_{(i+1)(j-5)} - a_{(i+1)(j-2)} - a_{(i+1)(j-3)}$ ,  $P_3 : a_{(i+2)(j)} - a_{(i+2)(j-1)} - a_{(i+2)(j-4)} - a_{(i+2)(j-6)} - a_{(i+2)(j-8)} - a_{(i+2)(j-7)} - a_{(i+2)(j-5)} - a_{(i+2)(j-2)} - a_{(i+2)(j-3)}$ ,  $P_4 : a_{(i)(j)} - a_{(i)(j-1)} - a_{(i)(j-4)} - a_{(i)(j-6)} - a_{(i)(j-8)}$  and  $\{E_3\} = \{e_1, e_2, e_3\}$ .

**Case (ii):** In  $G, d(a_{17}, a_{11}) = 3$  and the Hamiltonian path between the vertices  $a_{17}$  and  $a_{11}$  is  $P : \{P_1 \cup P_2 \cup P_3 \cup P_4\} \cup \{E_3\}$  where  $P_1 : a_{(i)(j-2)} - a_{(i)(j-5)} - a_{(i)(j-3)}$ ,  $P_2 : a_{(i+1)(j)} - a_{(i+1)(j-1)} - a_{(i+1)(j-4)} - a_{(i+1)(j-6)} - a_{(i+1)(j-8)} - a_{(i+1)(j-7)} - a_{(i+1)(j-5)} - a_{(i+1)(j-2)} - a_{(i+1)(j-3)}$ ,  $P_3 : a_{(i+2)(j)} - a_{(i+2)(j-1)} - a_{(i+2)(j-4)} - a_{(i+2)(j-6)} - a_{(i+2)(j-8)} - a_{(i+2)(j-7)} - a_{(i+2)(j-5)} - a_{(i+2)(j-2)} - a_{(i+2)(j-3)}$ ,  $P_4 : a_{(i)(j)} - a_{(i)(j-1)} - a_{(i)(j-4)} - a_{(i)(j-6)} - a_{(i)(j-7)} - a_{(i)(j-8)}$  and  $\{E_3\} = \{e_1, e_2, e_3\}$ .

**Case (iii):** Let  $t = 5$ . In  $G, d(a_{28}, a_{11}) = 5$  and the Hamiltonian path between the vertices  $a_{28}$  and  $a_{11}$  is  $P : \{P_1 \cup P_2 \cup P_3\} \cup \{E_2\}$  where  $P_1 : a_{(i+1)(j-1)} - a_{(i+1)(j)} - a_{(i+1)(j-4)} - a_{(i+1)(j-6)} - a_{(i+1)(j-8)} - a_{(i+1)(j-7)} - a_{(i+1)(j-5)} - a_{(i+1)(j-2)} - a_{(i+1)(j-3)}$ ,  $P_2 : a_{(i+2)(j)} - a_{(i+2)(j-1)} - a_{(i+2)(j-4)} - a_{(i+2)(j-6)} - a_{(i+2)(j-8)} - a_{(i+2)(j-7)} - a_{(i+2)(j-5)} - a_{(i+2)(j-2)} - a_{(i+2)(j-3)}$ ,  $P_3 : a_{(i)(j)} - a_{(i)(j-4)} - a_{(i)(j-1)} - a_{(i)(j-2)} - a_{(i)(j-3)} - a_{(i)(j-5)} - a_{(i)(j-7)} - a_{(i)(j-6)} - a_{(i)(j-8)}$  and  $\{E_2\} = \{e_2, e_3\}$ .

**Case (iv):** Let  $t = 7$ . In  $G, d(a_{22}, a_{11}) = 7$  and the Hamiltonian path between the vertices  $a_{22}$  and  $a_{11}$  is  $P : \{P_1 \cup P_2 \cup P_3\} \cup \{E_2\}$  where  $P_1 : a_{(i+1)(j-7)} - a_{(i+1)(j-8)} - a_{(i+1)(j-6)} - a_{(i+1)(j-4)} - a_{(i+1)(j)} - a_{(i+1)(j-1)} - a_{(i+1)(j-2)} - a_{(i+1)(j-5)} - a_{(i+1)(j-3)}$ ,  $P_2 : a_{(i+2)(j)} - a_{(i+2)(j-1)} - a_{(i+2)(j-4)} - a_{(i+2)(j-6)} - a_{(i+2)(j-8)} - a_{(i+2)(j-7)} - a_{(i+2)(j-5)} - a_{(i+2)(j-2)} - a_{(i+2)(j-3)}$ ,  $P_3 : a_{(i)(j)} - a_{(i)(j-4)} - a_{(i)(j-1)} - a_{(i)(j-2)} - a_{(i)(j-3)} - a_{(i)(j-5)} - a_{(i)(j-7)} - a_{(i)(j-6)} - a_{(i)(j-8)}$  and  $\{E_2\} = \{e_2, e_3\}$ .

From the above cases, it is obvious that the graph of  $H_3$  is  $(t_o^*)$  connected. □

**Lemma 2.2.** *The graph  $H_3$  is  $(t_e^*)$  connected for even  $t$  such that  $2 \leq t \leq \text{diam}(G)$ .*

*Proof.* Let  $G = H_n, n = 3$ . Clearly,  $\text{diam}(G) = 2^3 - 1 = 7$ . The vertex set and the edge set of  $H_3$  are same as in Lemma 2.1. We have the following cases. In each case, we take  $i = 1, j = 3^{n-1}$ .

**Case (i):** Let  $t = 2$ . In  $G$ ,  $d(a_{14}, a_{11}) = 2$  and the Hamiltonian path between the two vertices  $a_{14}$  and  $a_{11}$  is  $P : \{P_1 \cup P_2 \cup P_3 \cup P_4\} \cup \{E_3\}$  where  $P_1 : a_{(i)(j-5)} - a_{(i)(j-2)} - a_{(i)(j-3)}$ ,  $P_2 : a_{(i+1)(j)} - a_{(i+1)(j-1)} - a_{(i+1)(j-4)} - a_{(i+1)(j-6)} - a_{(i+1)(j-8)} - a_{(i+1)(j-7)} - a_{(i+1)(j-5)} - a_{(i+1)(j-2)} - a_{(i+1)(j-3)}$ ,  $P_3 : a_{(i+2)(j)} - a_{(i+2)(j-1)} - a_{(i+2)(j-4)} - a_{(i+2)(j-6)} - a_{(i+2)(j-8)} - a_{(i+2)(j-7)} - a_{(i+2)(j-5)} - a_{(i+2)(j-2)} - a_{(i+2)(j-3)}$ ,  $P_4 : a_{(i)(j)} - a_{(i)(j-1)} - a_{(i)(j-4)} - a_{(i)(j-6)} - a_{(i)(j-7)} - a_{(i)(j-8)}$  and  $\{E_3\} = \{e_1, e_2, e_3\}$ .

**Case (ii):** Let  $t = 4$ . In  $G$ ,  $d(a_{29}, a_{11}) = 4$  and the Hamiltonian path between the vertices  $a_{29}$  and  $a_{11}$  is  $P : \{P_1 \cup P_2 \cup P_3\} \cup \{E_2\}$  where  $P_1 : a_{(i+1)(j)} - a_{(i+1)(j-1)} - a_{(i+1)(j-4)} - a_{(i+1)(j-6)} - a_{(i+1)(j-8)} - a_{(i+1)(j-7)} - a_{(i+1)(j-5)} - a_{(i+1)(j-2)} - a_{(i+1)(j-3)}$ ,  $P_2 : a_{(i+2)(j)} - a_{(i+2)(j-1)} - a_{(i+2)(j-4)} - a_{(i+2)(j-6)} - a_{(i+2)(j-8)} - a_{(i+2)(j-7)} - a_{(i+2)(j-5)} - a_{(i+2)(j-2)} - a_{(i+2)(j-3)}$ ,  $P_3 : a_{(i)(j)} - a_{(i)(j-4)} - a_{(i)(j-1)} - a_{(i)(j-2)} - a_{(i)(j-3)} - a_{(i)(j-5)} - a_{(i)(j-7)} - a_{(i)(j-6)} - a_{(i)(j-8)}$  and  $\{E_2\} = \{e_2, e_3\}$ .

**Case (iii):** Let  $t = 6$ . In  $G$ ,  $d(a_{27}, a_{11}) = 6$  and the Hamiltonian path between the vertices  $a_{27}$  and  $a_{11}$  is  $P : \{P_1 \cup P_2 \cup P_3\} \cup \{E_2\}$  where  $P_1 : a_{(i+1)(j-2)} - a_{(i+1)(j-1)} - a_{(i+1)(j-4)} - a_{(i+1)(j-6)} - a_{(i+1)(j-8)} - a_{(i+1)(j-7)} - a_{(i+1)(j-5)} - a_{(i+1)(j-3)}$ ,  $P_2 : a_{(i+2)(j)} - a_{(i+2)(j-1)} - a_{(i+2)(j-4)} - a_{(i+2)(j-6)} - a_{(i+2)(j-8)} - a_{(i+2)(j-7)} - a_{(i+2)(j-5)} - a_{(i+2)(j-2)} - a_{(i+2)(j-3)}$ ,  $P_3 : a_{(i)(j)} - a_{(i)(j-4)} - a_{(i)(j-1)} - a_{(i)(j-2)} - a_{(i)(j-3)} - a_{(i)(j-5)} - a_{(i)(j-7)} - a_{(i)(j-6)} - a_{(i)(j-8)}$  and  $\{E_2\} = \{e_2, e_3\}$ .

From the above cases it is obvious that the graph of  $H_3$  is  $(t_e^*)$  connected.  $\square$

Lemma 2.1 and Lemma 2.2 leads to the following theorem.

**Theorem 2.3.** *The graph  $H_3$  is  $t^*$  connected for every  $t$  such that  $1 \leq t \leq \text{diam}(G)$ .*

**Lemma 2.4.** *The graph  $H_4$  is  $(t_o^*)$  connected for odd  $t$  such that  $1 \leq t \leq \text{diam}(G)$ .*

*Proof.* Let  $G = H_n$ ,  $n = 4$ . Clearly,  $\text{diam}(G) = 2^4 - 1 = 15$ . The graph of  $H_4$  contains three copies of  $H_3$  namely  $C_1$ ,  $C_2$  and  $C_3$  connected to each other by three unique edges  $\{E_{C_d}\}$  for  $d = 1, 2, 3$ . The edge  $\{E_{C_1}\}$  connects  $C_1$  and  $C_2$ ,  $\{E_{C_2}\}$  connects  $C_2$  and  $C_3$  and  $\{E_{C_3}\}$  connects  $C_3$  and  $C_1$ . Let the vertex set of  $H_4$  be  $V(H_4) = \{x_k, 1 \leq k \leq q\} \cup \{y_k, 1 \leq k \leq q\} \cup \{z_k, 1 \leq k \leq q\}$  and the edge set of  $H_4$  be  $E(H_4) = \{E_g, 1 \leq i \leq \frac{3^{n+1}-9}{6}\} \cup \{E'_g, 1 \leq i \leq \frac{3^{n+1}-9}{6}\} \cup \{E''_g, 1 \leq i \leq \frac{3^{n+1}-9}{6}\} \cup \{E_{C_d}\}$  where  $d = 1, 2, 3$  such that  $\{E_{C_1}\} = (x_p, y_q)$ ,  $\{E_{C_2}\} = (y_p, z_q)$ ,  $\{E_{C_3}\} = (z_p, x_q)$  for  $p = 3^{n-1} - m$ ,  $q = 3^{n-1}$ ,  $m = 2^{n-1} - 1$ ,  $r = s = 2^{n-2} - 1$ .

**Case (i):** For  $(x_i, x_j) \in C_1$ ,  $1 < i < q$ ,  $j = 1$ , the Hamiltonian path between  $x_i$  and  $x_j$  for  $d(x_i, x_j) = t$  where  $t = 1, 3, 5, 7$  is  $H_p : \{P_1 \cup P_2 \cup P_3 \cup P_4\} \cup \{E_{C_d}\}$ . Here  $P_1 : (x_i \rightarrow x_{p+r} \rightarrow x_p)$ ,  $P_2 : (y_q \rightarrow y_i \rightarrow y_p)$ ,  $P_3 : (z_q \rightarrow z_j \rightarrow y_p)$ ,  $P_4 : (x_q \rightarrow x_{q-s} \rightarrow x_j)$  and  $(E_{C_d})$  where  $d = 1, 2, 3$  such that  $\{E_{C_1}\} = (x_p, y_q)$ ,  $\{E_{C_2}\} = (y_p, z_q)$ ,  $\{E_{C_3}\} = (z_p, x_q)$ .

**Case (ii):** Let  $y_i \in C_2$  and  $x_j \in C_1$  where  $1 \leq i \leq q$ ,  $j = 1$ . Then the Hamiltonian path between  $(y_i, x_j)$  for  $d(y_i, x_j) = t$  where  $t = 9, 11, 13, 15$  is  $H_p : \{P_1 \cup P_2 \cup P_3\} \cup \{E_{C_d}\}$  where  $P_1 : (y_i \rightarrow y_{q-s} \rightarrow y_{p+r} \rightarrow y_p)$ ,  $P_2 : (z_q \rightarrow z_{q-s} \rightarrow z_{p+r} \rightarrow z_p)$ ,  $P_3 : (x_q \rightarrow x_{q-s} \rightarrow x_{p+r} \rightarrow x_j)$ , and  $(E_{C_d})$  where  $d = 2, 3$  such that  $\{E_{C_2}\} = (y_p, z_q)$ ,  $\{E_{C_3}\} = (z_p, x_q)$ .

Figures 2 and 3 depicts the Hamiltonian paths between  $x_i$  and  $x_j$  for  $d(x_i, x_j) = t$  where  $t = 1, 3, 5, 7$  and  $y_i$  and  $x_j$  for  $d(y_i, x_j) = t$  where  $t = 9, 11, 13, 15$ . □

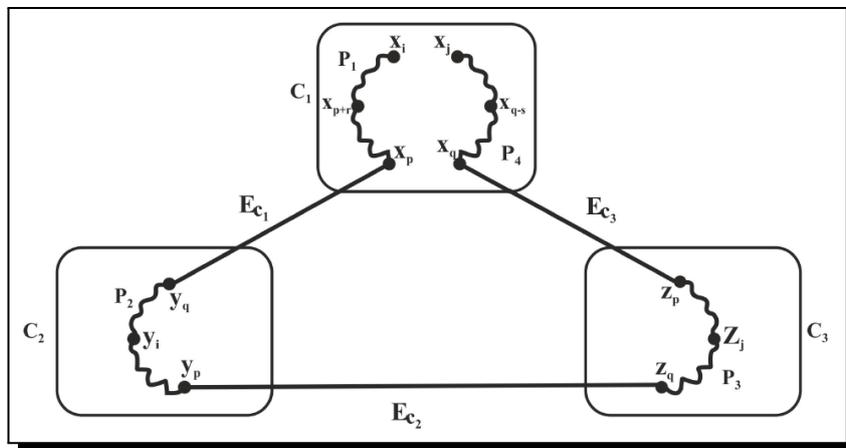


Figure 2

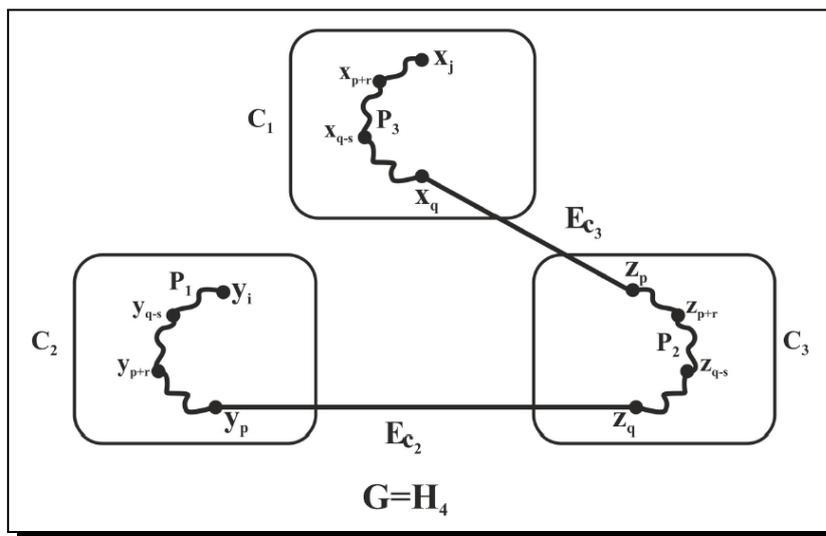


Figure 3

**Lemma 2.5.** The graph  $H_4$  is  $(t_e^*)$  connected for even  $t$  such that  $2 \leq t \leq \text{diam}(G)$ .

*Proof. Case (i):* Let  $(x_i; x_j) \in C_1$ ,  $1 < i < q$ ,  $j = 1$ . Then the Hamiltonian path between the two vertices  $x_i$  and  $x_j$  for  $d(x_i, x_j) = t$  where  $t = 2, 4, 6$  is  $P_1 : (x_i \rightarrow x_{p+r} \rightarrow x_p) \cup P_2 : (y_q \rightarrow y_i \rightarrow y_p) \cup P_3 : (z_q \rightarrow z_j \rightarrow y_p) \cup P_4 : (x_q \rightarrow x_{q-s} \rightarrow x_j) \cup (E_{C_d})$  where  $d = 1, 2, 3$  such that  $\{E_{C_1}\} = (x_p, y_q)$ ,  $\{E_{C_2}\} = (y_p, z_q)$ ,  $\{E_{C_3}\} = (z_p, x_q)$  for  $p = 3^{n-1} - m$ ,  $q = 3^{n-1}$ ,  $m = 2^{n-1} - 1$ ,  $r = s = 2^{n-2} - 1$ .

**Case (ii):** Let  $y_i \in C_2$  and  $x_j \in C_1$  where  $1 \leq i \leq q$ ,  $j = 1$ . Then the Hamiltonian path between  $(y_i, x_j)$  for  $d(y_i, x_j) = t$  where  $t = 8, 10, 12, 14$  is  $P_1 : (y_i \rightarrow y_{q-s} \rightarrow y_{p+r} \rightarrow y_p) \cup P_2 : (z_q \rightarrow z_{q-s} \rightarrow z_{p+r} \rightarrow z_p) \cup P_3 : (x_q \rightarrow x_{q-s} \rightarrow x_{p+r} \rightarrow x_j) \cup (E_{C_d})$  where  $d = 2, 3$  such that  $\{E_{C_2}\} = (y_p, z_q)$ ,

$\{E_{C_3}\} = (z_p, x_q)$  for  $p = 3^{n-1} - m$ ,  $q = 3^{n-1}$ ,  $m = 2^{n-1} - 1$ ,  $r = s = 2^{n-2} - 1$ . This proves that the graph of  $H_4$  is  $(t_e^*)$  connected.  $\square$

Lemma 2.4 and Lemma 2.5 leads to the following theorem.

**Theorem 2.6.** *The graph  $H_4$  is  $t^*$  connected for every  $t$  such that  $1 \leq t \leq \text{diam}(G)$ .*

**Theorem 2.7.** *The graph of  $H_n$  is  $t^*$  connected for all  $n > 4$  with  $d(u, v) = t$  such that  $1 \leq t \leq \text{diam}(G)$ .*

*Proof.* Let  $G = H_n$ ,  $n > 4$  be the Hanoi graph where  $n$  is a positive integer. Clearly,  $\text{diam}(G) = 2^n - 1$ . The graph of  $H_n$  contains three copies of  $H_{n-1}$  namely  $C_1$ ,  $C_2$  and  $C_3$  connected to each other by three unique edges  $\{E_{C_d}\}$  for  $d = 1, 2, 3$ . The edge  $\{E_{C_1}\}$  connects  $C_1$  and  $C_2$ ,  $\{E_{C_2}\}$  connects  $C_2$  and  $C_3$  and  $\{E_{C_3}\}$  connects  $C_3$  and  $C_1$ . The vertex set of  $H_n$  is  $V(H_n) = \{x_k \cup y_k \cup z_k\}$  where  $x_k \in C_1$ ,  $y_k \in C_2$ ,  $z_k \in C_3$  and  $1 \leq k \leq q$ ,  $q = 3^{n-1}$ . The edge set of  $H_n$  is  $E(H_n) = \left\{ E_{g/1} \mid 1 \leq g \leq \frac{3^{n+1}-9}{6} \cup E_{g'/1} \mid 1 \leq g' \leq \frac{3^{n+1}-9}{6} \cup E_{g''/1} \mid 1 \leq g'' \leq \frac{3^{n+1}-9}{6} \right\} \cup \{E_{C_d}\}$  where  $d = 1, 2, 3$  such that  $\{E_{C_1}\} = (x_p, y_q)$ ,  $\{E_{C_2}\} = (y_p, z_q)$ ,  $\{E_{C_3}\} = (z_p, x_q)$  for  $p = 3^{n-1} - m$ ,  $q = 3^{n-1}$ ,  $m = 2^{n-1} - 1$ ,  $r = s = 2^{n-2} - 1$ .

To establish the result, we consider the following cases.

**Case (i):** For the vertices  $(x_i, x_j) \in C_1$  where  $1 < i < q$ ,  $j = 1$  and  $d(x_i, x_j) = t$  for every odd  $t$ ,  $1 \leq t \leq \text{diam}(G)$ , the Hamiltonian path is  $P : \{P_1 \cup P_2 \cup P_3 \cup P_4\} \cup \{E_{C_d}\}$  as in Case (i) of Lemma 2.4.

**Case (ii):** For the vertices  $y_i \in C_2$  and  $x_j \in C_1$  where  $1 \leq i \leq q$ ,  $j = 1$  and  $d(x_j, y_i) = t$  for every odd  $t$ ,  $1 \leq t \leq \text{diam}(G)$ , the Hamiltonian path is  $P : \{P_1 \cup P_2 \cup P_3\} \cup \{E_{C_d}\}$  as in Case (ii) of Lemma 2.4.

**Case (iii):** For the vertices  $(x_i, x_j) \in C_1$  where  $1 < i < q$ ,  $j = 1$  and  $d(x_i, x_j) = t$  for every even  $t$ ,  $2 \leq t \leq \text{diam}(G)$ , the Hamiltonian path is  $P : \{P_1 \cup P_2 \cup P_3 \cup P_4\} \cup \{E_{C_d}\}$  as in Case (i) of Lemma 2.5.

**Case (iv):** For the vertices  $y_i \in C_2$  and  $x_j \in C_1$  where  $1 \leq i \leq q$ ,  $j = 1$  and  $d(x_j, y_i) = t$  for every even  $t$ ,  $2 \leq t \leq \text{diam}(G)$ , the Hamiltonian path is  $P : \{P_1 \cup P_2 \cup P_3\} \cup \{E_{C_d}\}$  as in Case (ii) of Lemma 2.5.  $\square$

### 3. Conclusion

In this paper we have shown that the Hanoi graphs  $H_n$  are Hamiltonian  $t^*$  connected for all  $n \geq 3$ . This concludes that the existence of Hamiltonian path in such networks (graphs) suffice to solve data communication problems.

### Acknowledgement

The first author thankfully acknowledges the support and encouragement provided by the Management, Director, Principal, Dean, HOD and Staff of BNM Institute of Technology,

Bengaluru. The authors are also thankful to the Management, HOD and R&D Centre, Department of Mathematics, Dr. Ambedkar Institute of Technology, Bengaluru.

### Competing Interests

The authors declare that they have no competing interests.

### Authors' Contributions

All the authors contributed significantly in writing this article. The authors read and approved the final manuscript.

## References

- [1] A. M. Hinz and D. Parisse, On the planarity of Hanoi graphs, *Expo. Math.* **20**(3) (2002), 263 – 268.
- [2] M. Annapoorna and R. Murali, Hamiltonian Laceability in the line graph of the  $(w, 1, n, k)$  graph, *International Journal of Computer Application* **5** (2015), 8 – 15.
- [3] C.-K. Li and I. Nelson, Perfect codes on the towers of Hanoi graph, *Bulletin of the Australian Mathematical Society* **57** (3) (1998), 367 – 376, DOI: 10.1017/s0004972700031774.
- [4] D. Arett and S. Doree, Colouring and counting on the tower of Hanoi graphs, *Mathematical Association of America* **83**(3) (2010), 200 – 209, DOI: 10.4169/002557010x494841.
- [5] D. Berend, A. Sapir and S. Solomon, The Tower of Hanoi problem on path graphs, *Discrete Applied Mathematics* **160** (2012), 1465 – 1483, DOI: 10.1016/j.dam.2012.02.007.
- [6] D. Berend and A. Sapir, The diameter of Hanoi graphs, *Information Processing Letters* **98**(2) (2006), 79 – 85, DOI: 10.1016/j.ipl.2005.12.004.
- [7] A. Girisha and R. Murali, Hamiltonian laceability in cyclic product and brick product of cycles, *International Journal of Graph Theory* **1**(1) (2013), 32 – 40.
- [8] A. M. Hinz, The tower of Hanoi, *Enseign. Math.* **35**(2) (1989), 289 – 321.
- [9] A. M. Hinz, Pascal's triangle and the tower of Hanoi, *Amer. Math. Monthly* **99** (1992), 538 – 544, DOI: 10.1080/00029890.1992.11995888.
- [10] A. M. Hinz, The tower of Hanoi, in: K. P. Shum, E. J. Taft and Z. X. Wan (eds.), *Algebras and Combinatorics*, Springer, Singapore, 277 – 289 (1999).
- [11] L. N. Shenoy and R. Murali, Laceability on a class of regular graphs, *International Journal of Computational Science and Mathematics* **2**(3) (2010), 397 – 406.
- [12] L. Xuemiao, Towers of Hanoi graphs, *International Journal of Computer Mathematics* **19**(1) (1986), 23 – 38.
- [13] R. Kumar and U. Maheswari, Matrix representation of Hanoi graphs, *International Journal of Science and Research* **4**(4) (2015), 173 – 174.
- [14] S. Aumann, K. A. M. Gotz, A. M. Hinz and C. Petr, The number of moves of the largest disc in shortest path on Hanoi graphs, *The Electronic Journal of Combinatorics* **21**(4) (2014), 1 – 22.
- [15] S. Klavzar and U. Milutinovic, Graphs  $S(n, k)$  and a variant of the tower of Hanoi problem, *Czechoslovak Math. J.* **47**(122) (1997), 95 – 104, URL: <http://dml.cz/dmlcz/127341>.